Output-Determinacy and Asynchronous Circuit Synthesis

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ACSD 2007



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Advantage of Decomposition

- Concurrency is typical for asynchronous circuits, but leads to state explosion
- Synthesis needs a representation of this state space
- Decomposition generates smaller components with smaller state spaces

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Determinism



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Dummy transitions are like λ/ϵ edges in an automaton. Inserted for convenience.



3 outputs are triggered by three inputs left: 6 places and 1 transitions – right: 9 places

If an action x^+ should occur after the action a^+ or b^+ occurs, one has to use either unsafe or non-deterministic STGs. However, unsafe STGs are not preferable in praxis.



Hiding of signals, i.e. labelling them with λ . Essential operation of our STG decomposition algorithm.



Hiding of a and b



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• Non-deterministic specifications can be **determinised**, i.e. transformed into a deterministic one with the same language.

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- For deterministic specifications the language is an adequate semantics

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- However, this leads to problems if the original specification is not output-determinate:

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- Non-deterministic specifications can be **determinised**, i.e. transformed into a deterministic one with the same language.
- For deterministic specifications the language is an adequate semantics
- However, this leads to problems if the original specification is not output-determinate:
 - Violation of semi-modularity
 - Violation of correctness

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Output-Determinacy – Why not Determinise?

Violation of Semi-Modularity



left: non-output-determinate: STG right: determinised version, x^+ is in conflict with y^+



Violation of Correctness



left: non-output-determinate: STG right: determinised version, can be synthesised but is not correct

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- Non-output-determinacy is also a CSC conflict
- CSC conflicts can be resolved by insertion of internal signals in a behaviour-preserving way
- Does not work for violation of output-determinacy

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A non-output-determinate specification is illformed and cannot be implemented speedindependently.

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A non-output-determinate specification is illformed and cannot be implemented speedindependently.

The semantics of an output-determinate speed-independent specification is its language.

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These give a semantics for non-deterministic specifications.

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 \bullet Decomposition is not always so easy: often, it is impossible to remove all λ transitions introduced before

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- \bullet Decomposition is not always so easy: often, it is impossible to remove all λ transitions introduced before
- Then: **backtracking**, i.e. restarting at the initial component with additional signals

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- \bullet Decomposition is not always so easy: often, it is impossible to remove all λ transitions introduced before
- Then: **backtracking**, i.e. restarting at the initial component with additional signals
- Either necessary to get a correct component, or just for technical reasons

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- \bullet Decomposition is not always so easy: often, it is impossible to remove all λ transitions introduced before
- Then: **backtracking**, i.e. restarting at the initial component with additional signals
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New Approach

- No backtracking for technical reasons
- λ transitions can be left in the component, as long as it is output-determinate (checked implicitly during synthesis)

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Decomposition – Non-Secure Contraction



Language is changed

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Decomposition – Non Safeness-Preserving Contraction





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Decomposition – Non Safeness-Preserving Contraction



- Synthesis tools do not support non-safe STGs
- or just with a place capacity
- Latter is inefficient: every contraction doubles potential place capacity

Benchmarks — SEQPARTREE



Artificial BALSA Handshake component tree of sequencers and parallelisers with 3 Levels

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| Level | P - T | In - Out | Old | New | Synthesis |
|-------|--------------|-------------|-------|------|-----------|
| 5 | 382 – 252 | 33 – 93 | 1 | 1 | 5 |
| 6 | 798 – 508 | 65 - 189 | 4 | 4 | 16 |
| 7 | 1566 - 1020 | 129 - 381 | 9 | 8 | 22 |
| 8 | 3230 - 2044 | 257 – 765 | 32 | 17 | 1:02 |
| 9 | 6302 - 4092 | 513 - 1533 | 1:27 | 1:18 | 1:30 |
| 10 | 12958 - 8188 | 1025 - 3069 | 42:37 | 6:03 | 4:32 |

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Synthesis of an STG with 4094 signals in about 11 minutes.

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Deciding Output-Determinacy ...

- $\bullet \ \mathrm{PSPACE}$ complete for bounded/safe STGs
- $\bullet \ \mathrm{ExpSpace}$ hard for unbounded STGs

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Deciding Output-Determinacy ...

- $\bullet \ \mathrm{PSPACE}$ complete for bounded/safe STGs
- $\bullet \ \mathrm{ExpSpace}$ hard for unbounded STGs

Decomposition:

- Simplified correctness notion
- Decomposition can be applied to output-determinate (possibly) non-deterministic specifications
- More decomposition operations

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Conclusion

Introduced output-determinacy

- Relaxation of determinism and determinacy
- Fits very well with STGs and speed-independency
- Allows a language-based semantics

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Introduced output-determinacy

- Relaxation of determinism and determinacy
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- Allows a language-based semantics

Applied output-determinacy to STG decomposition

- Decomposition gets faster
- Allows more STG specifications
- More decomposition operations

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Future Reseach

- Semantics for internal signals application to output-determinacy and decomposition
- Decomposition works structurally only.
 Fast but problems with 'tricky' specifications
- \bullet Apply decomposition to BALSA resynthesis
- Preserving CSC in the components

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Future Reseach

- Semantics for internal signals application to output-determinacy and decomposition
- Decomposition works structurally only.
 Fast but problems with 'tricky' specifications
- \bullet Apply decomposition to BALSA resynthesis
- Preserving CSC in the components

Related Work

- Carmona and Cortadella: ILP approach, e.g ILP models for the synthesis of asynchronous control circuits
- Myers and Yoneda: structural decomposition, e.g. Synthesis of Speed Independent Circuits Based on Decomposition